Continuation and Exceptions

Control Flow In Sequential Languages

Imperative Programming Control Flow of Programs

Structured control flow

- Sequence of statements
 a:= b; b := c; }
- Conditional
 - if (a < b) then c else d;
 - switch(a){...}
- Loops
 - □ for (...) {...};
 - while (...) {...};
- Jumping out of a block
 - break, continue, return,...
- Non-structured control flow
 - Goto, conditional jump
 - Used to implement structured control flow in assembly

Controlling Jumps

Structured jumps

if ... then ... else ... end
while ... do ... end
for ... { ... }
case ...

- Group code in logical blocks
- Avoid uncontrolled jumps, e.g., into the middle of a block
- Focus of this chapter: quickly jumping into and out of a program in an organized fashion
 - Jumping right into the mid of a block ---- continuation passing
 - The scenario: my task was interrupted, now I want to resume from where I stopped
 - Jumping out from the mid of a block? ---- exception handling
 - The scenario: something unexpected happened; need to jump out until some caller knows what to do with the errors.

Continuations

• Capture the continuation at some point to be used later

A function (closure) that takes a single parameter, the result of the past evaluation, and returns the result of the entire program.

Save the entire runtime environment as a closure

- Code pointer: where to start evaluating the instructions
- Environment pointer: the entire relevant memory stores
- To jump into the mid of a program, make a function call to the continuation
- Useful in
 - Implementing functional programming languages
 - Operating system scheduling, Web site design
 - The scenario: my task was interrupted, now I want to resume from where I stopped

Continuation of Expressions

Continuation: impose sequential ordering in sub-expressions

- The continuation of an expression is "the remaining work to be done after evaluating the expression"
- Continuation of e is a function applied to the result of e
- Enforce evaluation order in functional languages
 - Evaluate current expression
 - Save the result into a variable
 - Evaluate the rest of the computation

```
2*x + 3*y + 1/x + 2/y
let val r2x = 2 * x in

let val r3y = 3 * y in

let val sum1=r2x + r3y in

let val r1x = 1 / x in

let val sum2 = sum1 + r1x in

let val r2y = 2 / y in

sum2 + r2y

end

.....end
```

let
$$r2x = 2*x$$
 in ... end
is equivalent to
 $(fn r2x=> ...) (2 * x)$

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Continuation and Tail Calls

• A function call from g to f is a tail call

- if g returns the result of calling f with no further computation
- Example (red: tail call; blue: non-tail call) fun f(x) = if x > 0 then x else f(x+1)*2 fun f(x,y) = if x>y then x else f(2*x,y);
- Tail calls do not need to return to caller
 - Can we convert all functions to tail recursion?
 - If a program needs to be re-enterable, function calls shouldn't return to caller
 - Solution: continuation passing
 - Pass continuation as parameter to callee
 - Callee does not need to return to caller

Continuation Passing

Standard function

 fun fact(n) = if n=0 then 1 else
 n*fact(n-1)

 Continuation form

 fun fact(n, K) =
 if n=0 then K(1)
 else fact(n-1,fn x=>K(n*x));
 fact(n, fn x=>x) computes n!

 Example computation

 fun fact(n, fn x=>x)

fact(3,fn x=>x)= fact(2,fn y=>((fn x=>x)(3*y))) = fact(1, fn x=>((fn y=>3*y)(2*x))) = fn x=>((fn y=>3*y)(2*x)) 1 = 6 **•** For each function definition F

- Extend the definition with a continuation parameter K
- At each function call inside F
 - Convert the rest of computation into a new continuation function
 - Convert f into a tail call, which takes the new continuation function as an extra argument.
- At each normal return
 - Return the result of invoking continuation K with the original returned value

General uses of continuations

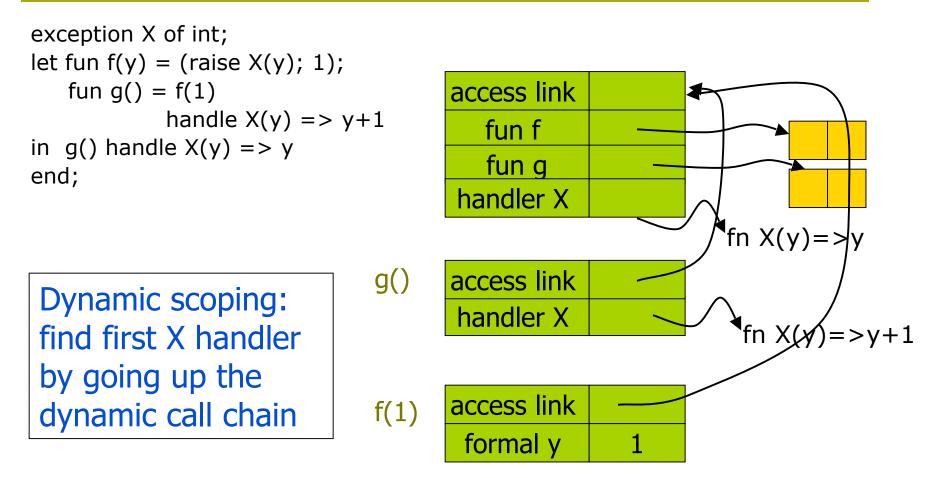
- Explicit control
 - Normal termination -- call continuation
 - Abnormal termination -- do something else
- Compilation techniques
 - Call to continuation is functional form of "go to"
 - Jump to the middle of a block by saving the environment in the function closure and restore the environment before jump
- Web applications, Web Services, MOM and SOA services
 - Handle long running workflows
 - Workflow may take 1 year to complete
 - Progress of subtasks is asynchronous
 - Sequential programming is simpler than asynchronous
- Continuations provide
 - An easy way to suspend workflow execution at a wait state
 - Thread of control can be resumed when the next message/event occurs, maybe some long time ahead

Exception: Structured Exit

□ When something unusual happens, we want a program to

- Jump out of one or many levels of nested blocks
- Until reaching some program point to continue
- Pass information to the continuation point
- May need to free heap space, other resources
- An exception is a dynamic jump
 - Don't know where to resume execution until runtime
 - Jump out of current block
 - Look for a matching exception handler in most recently entered blocks
 - General dynamic scoping rule
 - Multiple functions could handle the same exception
 - Jump to most recently established handler on run-time stack
 - Callers know how to handler error, defining block doesn't

Dynamic Scoping of Handler



When Should We Use Exceptions?

Separation of concern: handle unusual situations

- Examples: division by zero, null pointers, unexpected inputs
- When exceptions are handled, error recovery
- Otherwise, evaluation aborts on error conditions
- Flexible control flow
 - Return immediately to where the error can be handled
 - Jump out multiple blocks at a time
- What languages have exception support?
 - C++, Java, ML, Ada, ...

Defining Exceptions

Exception declaration

- Type of data that can be passed in exception
 - ML: exception <name> of <type>
 - C++/Java: any data type
- Raising an exception
 - Abort the rest of current block and jump out
 - ML: raise <name> <arguments>;
 - C++: throw <value>;
- Handling an exception
 - Continue normal execution after exception
 - ML: <exp1> handle <pattern>=><exp2>; ...
 - □ C++: try { ...} catch (<type> var) {...} ...

Exceptions vs. Type System

Are exceptions part of the type system?

- Raising expressions: not part of the type system
 - Expression e has type t if normal termination of e produces value of type t
 - Raising exception is not normal termination
 - Example: 1 + raise X is not valid
- Handling exceptions (=> (value))
 - Converts exception to normal termination
 - Need type agreement
 - 1 + ((raise X) handle X => e) Type of e must be int
 - 1 + (e_1 handle X => e_2) Type of e_1 , e_2 must be int

How Are Exceptions Handled?

ML

- exception X of int; let fun f(y) = (raise X(y); 1); fun g() = f(1) handle X(y) => y+1 in g() handle X(y) => y end;
- What are the events that have occurred?
 - Enter the let expression
 - Make function call g()
 - Make function call f(1)
 - Function call f(1) raises exception X(1)
 - Exception X(1) is handled in function call g()
 - Function call g() returns with value 2
 - The let expression exits

Exception vs. Continuation

Continuation

- Explicitly represent the rest of computation
- Do not need to return to the caller
 - Can use exception to avoid returning to the caller
- Raising exception
 - Jumping out of multiple blocks at a time
 - Different continuation for normal and exceptional situations

Continuation of exception

- rest of computation after exception is handled
- Raising exceptions may have complications
 - Resource management: opened files, garbage collection
 - Use continuation passing to implement exception
 - Pass multiple continuations: one to handle normal condition, the others to handle exceptions